

Topological closure of ideals of commutative formal power series and applications

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I. Introduction

- ▷ Rewriting on topological structures
- ▷ Formal power series as Cauchy-completion of polynomials

II. Ideals of formal power series

- ▷ Standard bases and topological confluence
- ▷ Topological closure of ideals

III. Applications

- ▷ Equivalence of confluences for formal power series
- ▷ Relation between standard bases and topological confluence

IV. Conclusion and perspectives

I. INTRODUCTION

Rewriting theory

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- Involution divisions

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- Involutive divisions

Abstraction

Abstract rewriting theory

Abstract properties common to all concrete rewriting systems:
termination, **confluence**, **normal forms**

Abstract Rewriting System

→ A an underlying set

→ \rightarrow a binary relation on A

We write $a \rightarrow b$ for $(a, b) \in \rightarrow$

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 $a = a_0 \rightarrow a_1 \rightarrow \dots \rightarrow a_\ell = b$

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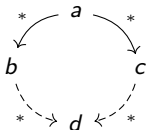
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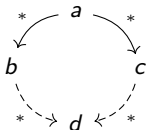
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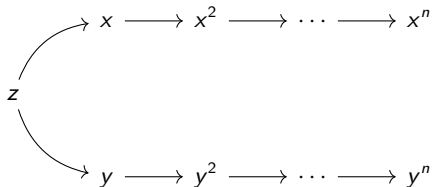
Confluence**Example**

Multivariate division with respect to R is **confluent** iff R is a **Gröbner basis**

Confluence “at the limit”

In $\mathbb{K}[[x, y, z]]$ with the inverse deglex order such that $z > y > x$ take

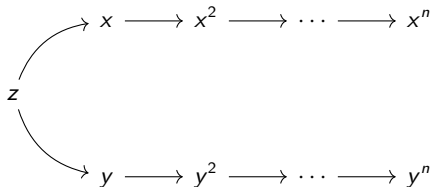
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$$R = \{z - y, z - x, y - y^2, x - x^2\}.$$



The two branches will never have a common element

Hence the system is **not confluent**

However with the (x, y, z) -adic topology both branches converge to 0

Topological Abstract Rewriting System

→ (X, τ) a **topological space**

→ \rightarrow a binary relation on X

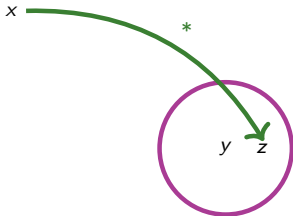
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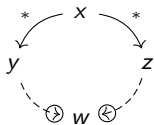
Topological rewriting relation

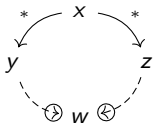
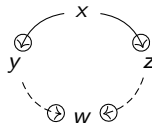
Write $x \dashrightarrow y$ if for **every neighbourhood** U of y **there exists** $z \in U$ s.t. $x \xrightarrow{*} z$

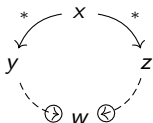
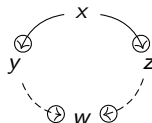


Note how $x \xrightarrow{*} y$ implies $x \dashrightarrow y$

Topological confluence



Topological confluence**Infinitary confluence**

Topological confluence**Infinitary confluence****Strength of confluences**

For every TARS we have:

confluence \implies **topological confluence**

infinitary confluence \implies **topological confluence**

Discrete rewriting system

If $x \rightarrow_{\oplus} y$ implies $x \xrightarrow{*} y$, then we say that the TARS (X, τ, \rightarrow) has **discrete rewriting**.

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In such a case, **confluence**, **topological confluence** and **infinitary confluence** are **trivially equivalent**.

For instance, if τ is the **discrete topology**, then (X, τ, \rightarrow) has **discrete rewriting**.

Line with two origins

$$X := (\mathbb{R} \times \{\pm 1\}) / \sim$$

where $(x, 1) \sim (x, -1)$ if $x \neq 0$

$$\forall n \in \mathbb{N}, \quad \left(\frac{1}{2^n}, 1\right) \rightarrow \left(\frac{1}{2^{n+1}}, 1\right)$$

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$$\left(\frac{1}{2}, 1\right) = \left(\frac{1}{2}, -1\right)$$



$$\left(\frac{1}{4}, 1\right) = \left(\frac{1}{4}, -1\right)$$

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Cyclic relation

$$X := [0, 2] \subset \mathbb{R}$$

$$\frac{1}{2^{n+1}} \begin{array}{c} \curvearrowright \\ \curvearrowleft \end{array} \frac{1}{2^n} \quad 2 - \frac{1}{2^n} \begin{array}{c} \curvearrowright \\ \curvearrowleft \end{array} 2 - \frac{1}{2^{n+1}}$$

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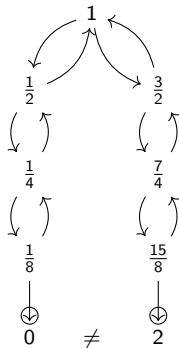
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Note how $(n, m) \xrightarrow{*} (n', m')$ iff $n \leq n'$ and $m \leq m'$

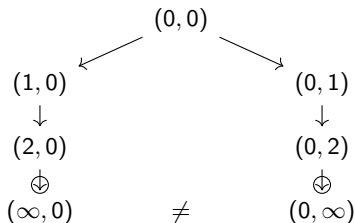
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$$\text{val}(xy^2z^2 + z^3 + y) = 1$$

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**Metric**

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**Cauchy-completion**

Cauchy sequences in $\mathbb{K}[x_1, \dots, x_n]$
modulo $(x_n)_n \sim (y_n)_n$ if and only
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Formal power series

Denoted by $\mathbb{K}[[x_1, \dots, x_n]]$

Structure on $\mathbb{K}[[x_1, \dots, x_n]]$

Algebra operations on $\mathbb{K}[x_1, \dots, x_n]$ are **continuous**

Hence **extend naturally** on $\mathbb{K}[[x_1, \dots, x_n]]$

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Formal power series as infinite linear combinations of monomials

$\mathbb{K}[[x_1, \dots, x_n]]$ isomorphic to the **dual vector space** of $\mathbb{K}[x_1, \dots, x_n]$

Since the monoid $[x_1, \dots, x_n]$ of monomials is a basis of $\mathbb{K}[x_1, \dots, x_n]$

$$f \in \mathbb{K}[[x_1, \dots, x_n]] \quad \leftrightarrow \quad f : [x_1, \dots, x_n] \rightarrow \mathbb{K} \text{ a map}$$

Denote $\langle f | m \rangle := f(m)$ then $f =: \sum_{m \in [x_1, \dots, x_n]} \langle f | m \rangle m$

Theorem. [Chenavier, Cluzeau, ML, 2024]

Let I be an ideal of commutative formal power series.

Given any f in the topological closure of I , we can compute a cofactor representation of f with respect to a system of generators of I .

In other words, we prove constructively that I is **topologically closed**.

Theorem. [Chenavier, Cluzeau, ML, 2024]

Let R be a set of commutative formal power series and $<$ be a local monomial order that is compatible with the degree.

The rewriting system induced by R and $<$ is **topologically confluent** if and only if it is **infinitary confluent**.

II. IDEALS OF FORMAL POWER SERIES

Local monomial orders

→ Total order **compatible** with monomial multiplication

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Rewriting on formal power series: **same as multivariate division on polynomials** but with respect to

- a **local** order $<$ **compatible with the degree**
- a set R of non-zero formal power series

Standard bases

Defined syntactically like **Gröbner bases** for polynomials *i.e.*
a subset $G \subseteq I$ of an ideal

$$\langle \text{LM}(G) \rangle = \text{LM}(I)$$

(Note: LM is w.r.t. the local order)

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Theorem [Chenavier, 2020].

R is a **standard basis** of the ideal it generates for a **local** order $<$ if, and only if, the system induced by R and $<$ is **topologically confluent**.

Counter-example of topological confluence \Rightarrow confluence

Consider again, in $\mathbb{K}[[x, y, z]]$

$$R = \{z - y, z - x, y - y^2, x - x^2\}.$$

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- $\rightarrow \text{LM}(R) = \{x, y, z\}$ and
- \rightarrow if $f \in I(R)$ then f has no constant coefficient

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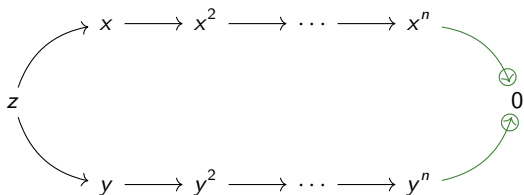
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Thus the system is **topologically confluent**



However we saw previously that it is **not confluent**

Ideals of formal power series are topologically closed

→ $\mathbb{K}[[x_1, \dots, x_n]]$: local noetherian topological ring with respect to the (x_1, \dots, x_n) -adic topology. Therefore a **Zariski ring**
[Samuel, Zariski, 1975]

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[Samuel, Zariski, 1975]
- Constructive proof providing a **cofactor representation** of a formal power series in the topological closure of the ideal
[Chenavier, Cluzeau, ML, 2024]

Lemma. Let I be an ideal in $\mathbb{K}[[x_1, \dots, x_n]]$ and $<$ be a **local** monomial order **compatible with the degree**

Then $\text{LM}(\bar{I}) = \text{LM}(I)$ where \bar{I} denotes the **topological closure of I**

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Proof. If $f \in \bar{I} \setminus \{0\}$ then $\exists (f_k)_k$ in I converging to f
Take f_k such that $\delta(f_k, f) < \frac{1}{2^{\deg(\text{LM}(f))}}$

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Thus $\deg(\text{LM}(f_k - f)) > \deg(\text{LM}(f))$

By **compatibility with the degree** we get $\text{LM}(f_k - f) < \text{LM}(f)$

This means that for all $m \geq \text{LM}(f)$ we have $\langle f_k | m \rangle = \langle f | m \rangle$

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This means that for all $m \geq \text{LM}(f)$ we have $\langle f_k | m \rangle = \langle f | m \rangle$

Hence $\text{LM}(f) = \text{LM}(f_k)$ but $f_k \in I$ therefore $\text{LM}(f) \in \text{LM}(I)$

Goal of the proof

Let I be an ideal in $\mathbb{K}[[x_1, \dots, x_n]]$

Let $<$ a **local** monomial order **compatible with the degree**

Fix $G := \{s_1, \dots, s_\ell\}$ a **standard basis** of I with respect to $<$

Let f be **in the topological closure** of I

Construct $(f_1, \dots, f_\ell) \in \mathbb{K}[[x_1, \dots, x_n]]^\ell$ such that $f = f_1 s_1 + \dots + f_\ell s_\ell$.

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Strategy

Construct a sequence $(f_i^{(k)})_{k \in \mathbb{N}}$ for each $i \in [1.. \ell]$

Prove that they are Cauchy

Take their limits and show that they yield a cofactor representation of f

Consider

$$F_k := f - \sum_{i=1}^{\ell} f_i^{(k)} s_i \in \bar{I}$$

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Otherwise, we can eliminate $m_k := \text{LM}(F_k)$ with the **standard basis** G by choosing an $i_k \in \llbracket 1 .. \ell \rrbracket$ and a $q_k \in [x_1, \dots, x_n]$ such that

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We define $f_i^{(k+1)} := f_i^{(k)}$ for all $i \neq i_k$ and

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By induction hypothesis it follows that:

$$F_{k+1} := \sum_{i=1}^{\ell} f_i^{(k+1)} s_i \in \bar{I}$$

Facts

Following from the facts that:

- we have **finitely many** variables
- $<$ is **compatible with the degree**
- the sequence $(m_k)_k$ of eliminated monomials is strictly decreasing

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Following from the facts that:

- we have **finitely many** variables
- $<$ is **compatible with the degree**
- the sequence $(m_k)_k$ of eliminated monomials is strictly decreasing

we have that the sequences $(f_i^{(k)})_k$ are Cauchy for any $i \in \llbracket 1 .. \ell \rrbracket$

Denote by $f_i^{(\infty)}$ their respective limits

Then by continuity of the metric and the facts above it follows that:

$$\delta \left(f - \sum_{i=1}^{\ell} f_i^{(\infty)} s_i, 0 \right) = 0$$

hence our desired result

III. APPLICATIONS

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Theorem. [Chenavier, Cluzeau, ML, 2024]

Let R be a set of formal power series and $<$ be a **local** monomial order that is **compatible with the degree**.

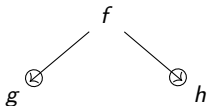
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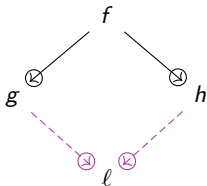


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Close the diagram

- Fix R a non-empty set of non-zero formal power series
- Fix $<$ a **local** monomial order **compatible with the degree**
- Write \rightarrow the one-step rewriting relation induced by R and $<$

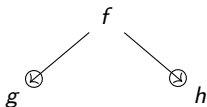
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Let $f, g, h \in \mathbb{K}[[x_1, \dots, x_n]]$ such that:



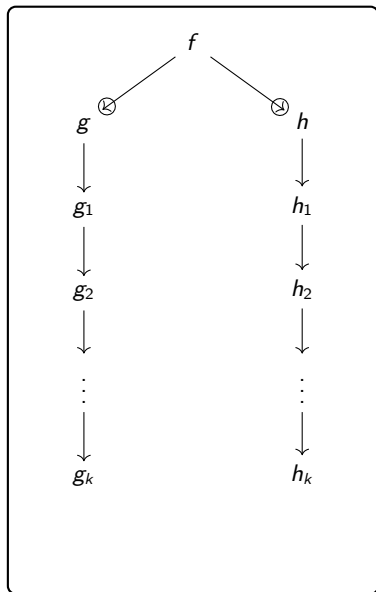
Goal

Construct inductively **two rewriting sequences** starting from g and h respectively that will be proven to be **Cauchy**

It will turn out that the limits are then equal and hence give a **common topological successor** to g and h

→ By induction:

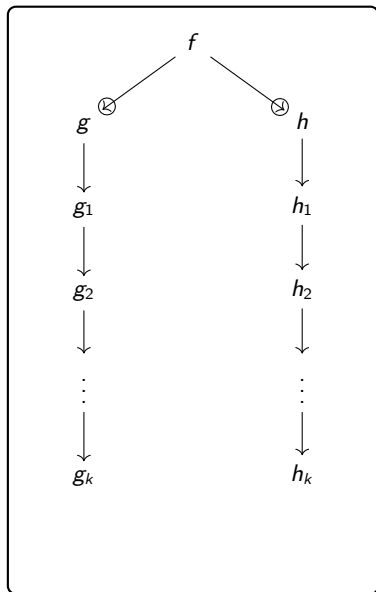
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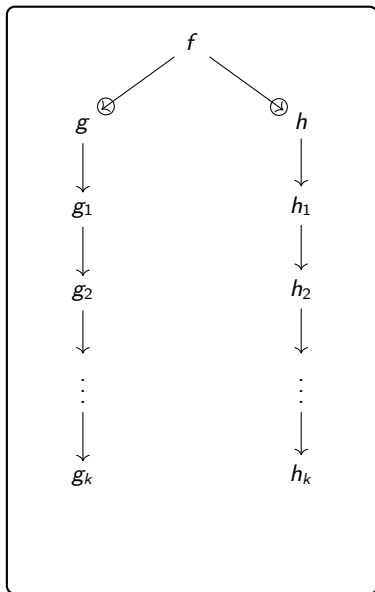
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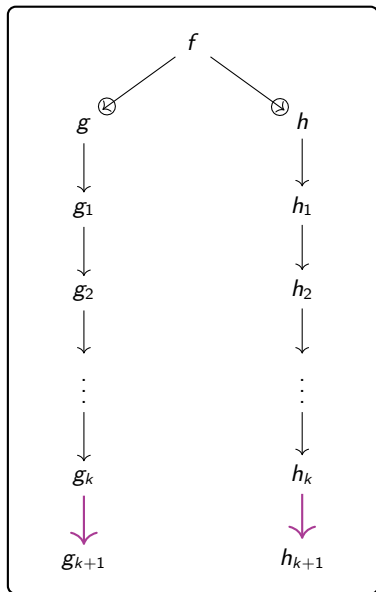
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→ Rewrite LM $(g_k - h_k)$



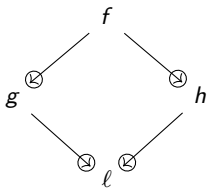
Facts

- the sequences $(g_k)_{k \in \mathbb{N}}$ and $(h_k)_{k \in \mathbb{N}}$ are **Cauchy**
- their limits are **equal**

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- their limits are **equal**

So $\lim_{k \rightarrow \infty} g_k = \lim_{k \rightarrow \infty} h_k =: \ell$



Which shows that \rightarrow is **infinitary confluent**

Let R be a set of non-zero formal power series.

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Write \rightarrow the relation induced by R and $<$ and \rightarrow_{\oplus} the topological rewriting relation associated to it.

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Write \rightarrow the relation induced by R and $<$ and \rightarrow_{\oplus} the topological rewriting relation associated to it.

Then the following properties are equivalent:

- (i) the system is **topologically confluent**
- (ii) for all $f \in I$, we have $f \rightarrow_{\oplus} 0$
- (iii) for all $f \in I \setminus \{0\}$, we have f reducible
- (iv) for all $f \in I \setminus \{0\}$, we have $\text{LM}(f)$ reducible
- (v) R is a **standard basis**
- (vi) the set of normal forms forms a canonical set of representatives for the quotient algebra $\mathbb{K}[[x_1, \dots, x_n]]$ modulo $I(R)$

IV. CONCLUSION AND PERSPECTIVES

Conclusion and perspectives

Summary of presented notions and results:

- ▷ we introduced the basic ideas of **topological rewriting theory**
- ▷ we proved constructively that ideals of commutative formal power series are **topologically closed**
- ▷ we showed that **topological confluence** is equivalent to **infinitary confluence** for formal power series

Further works:

- ▷ show that the topological rewriting relation induces convergent rewriting chains in the context of formal power series
- ▷ adapt Newman's lemma to topological rewriting theory
- ▷ develop computational tools to study Taylor series and formal solutions of PDEs

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THANK YOU FOR LISTENING!